

Pointless Tainting?

Evaluating the Practicality of Pointer
Tainting

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Why pointer tainting?

- Attacks

Exploit low-level memory errors

Buffer overflows

Dangling pointers

Format strings

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- Keyloggers, etc.

Installed by users or by the way of exploits

e.g., trojan

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- **Keyloggers, etc.**

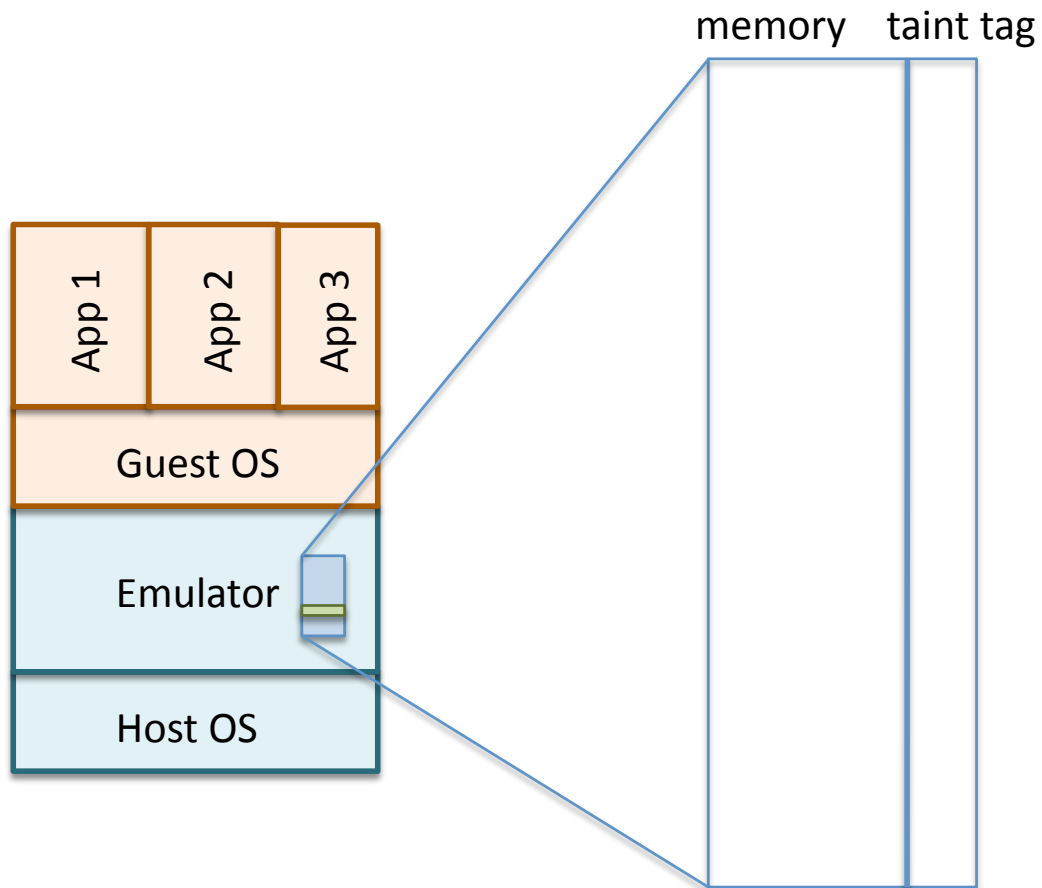
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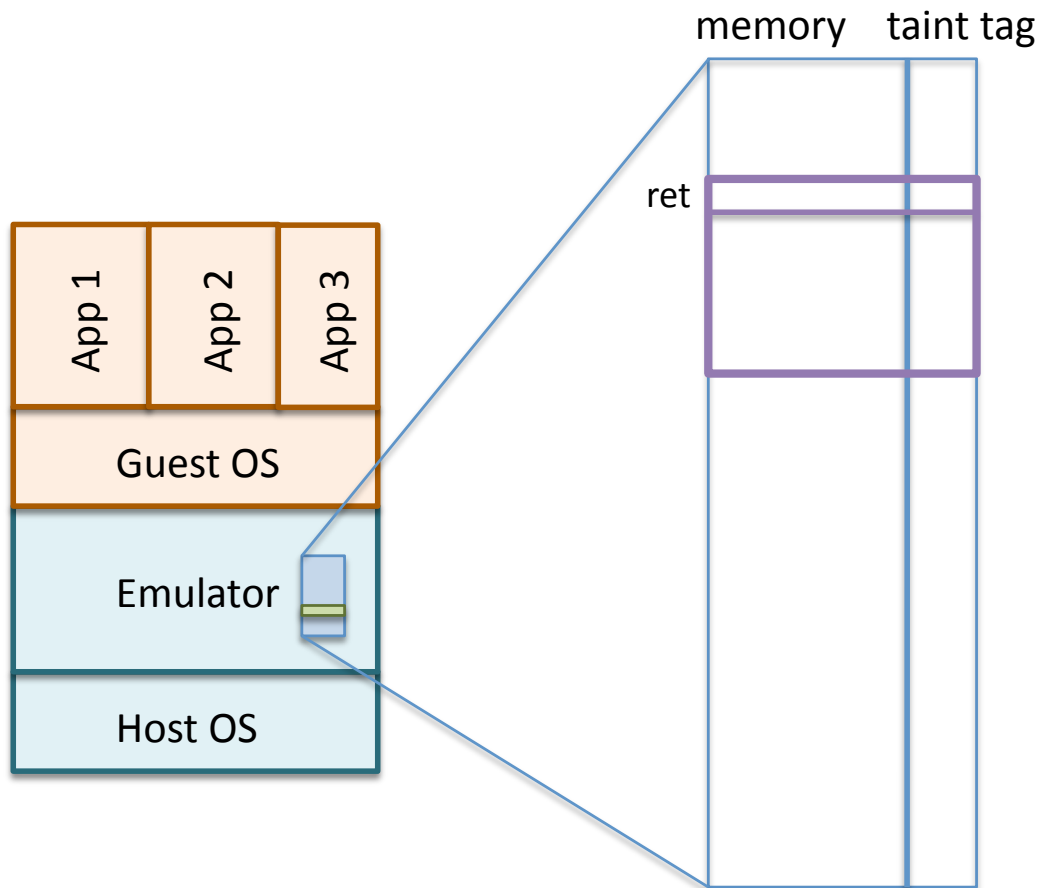
- **Pointer tainting**

- Capable of detecting
 - Memory corruption attacks
 - Both control- and non-control-diverting
 - Privacy-breaching malware
- PROBLEMATIC

Basic tainting

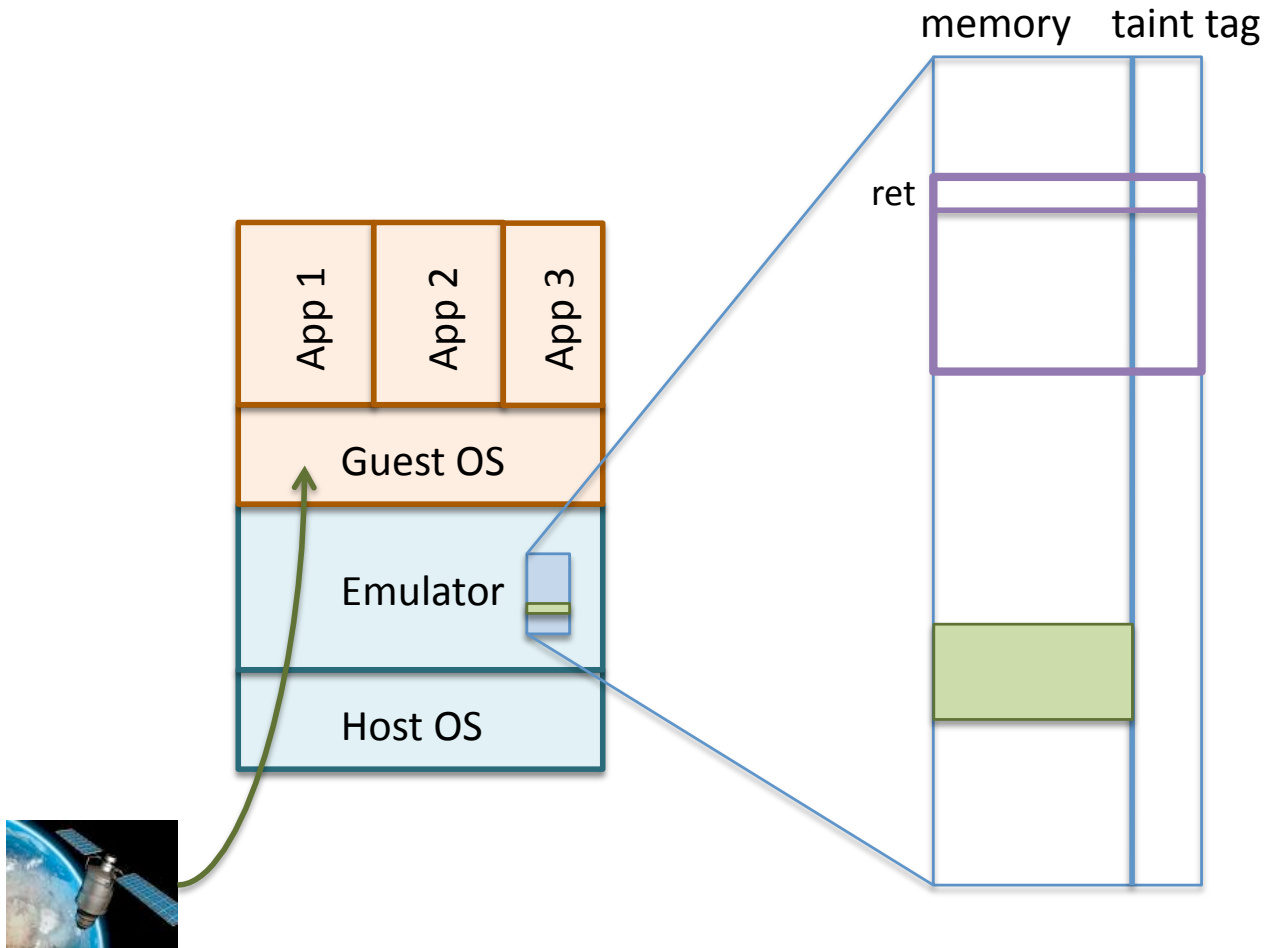


Basic tainting



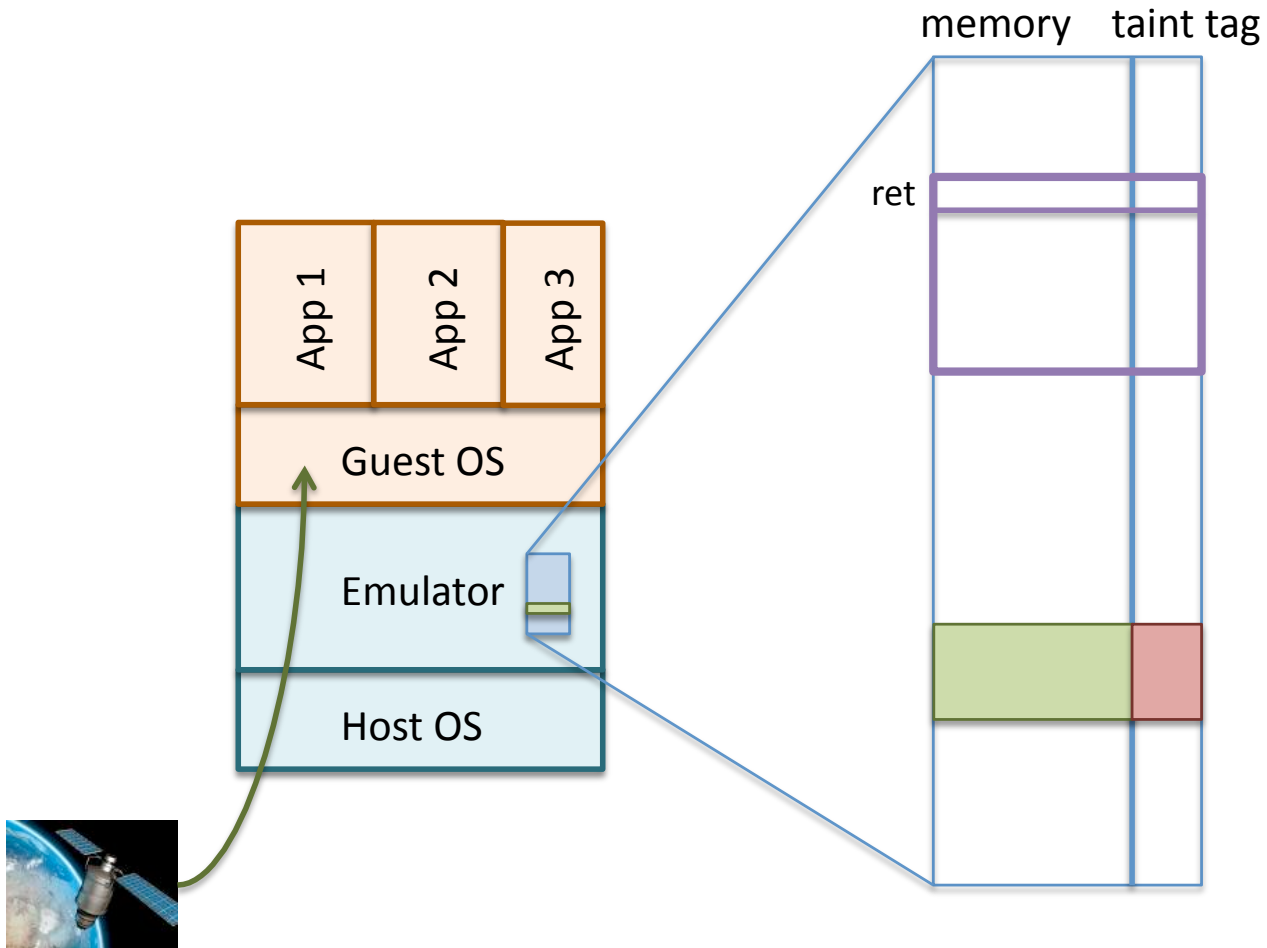
Basic tainting

1. Mark network data as tainted.



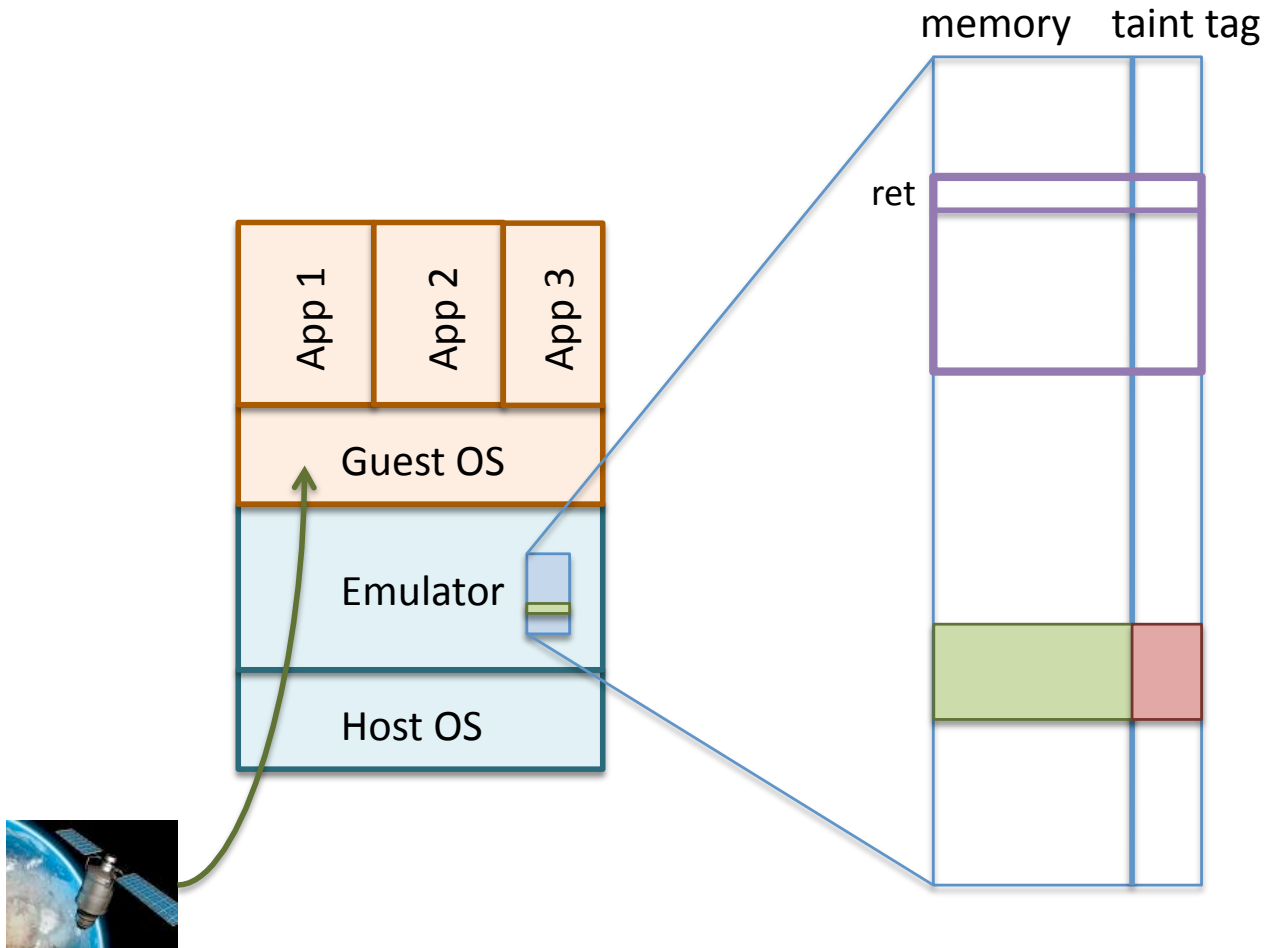
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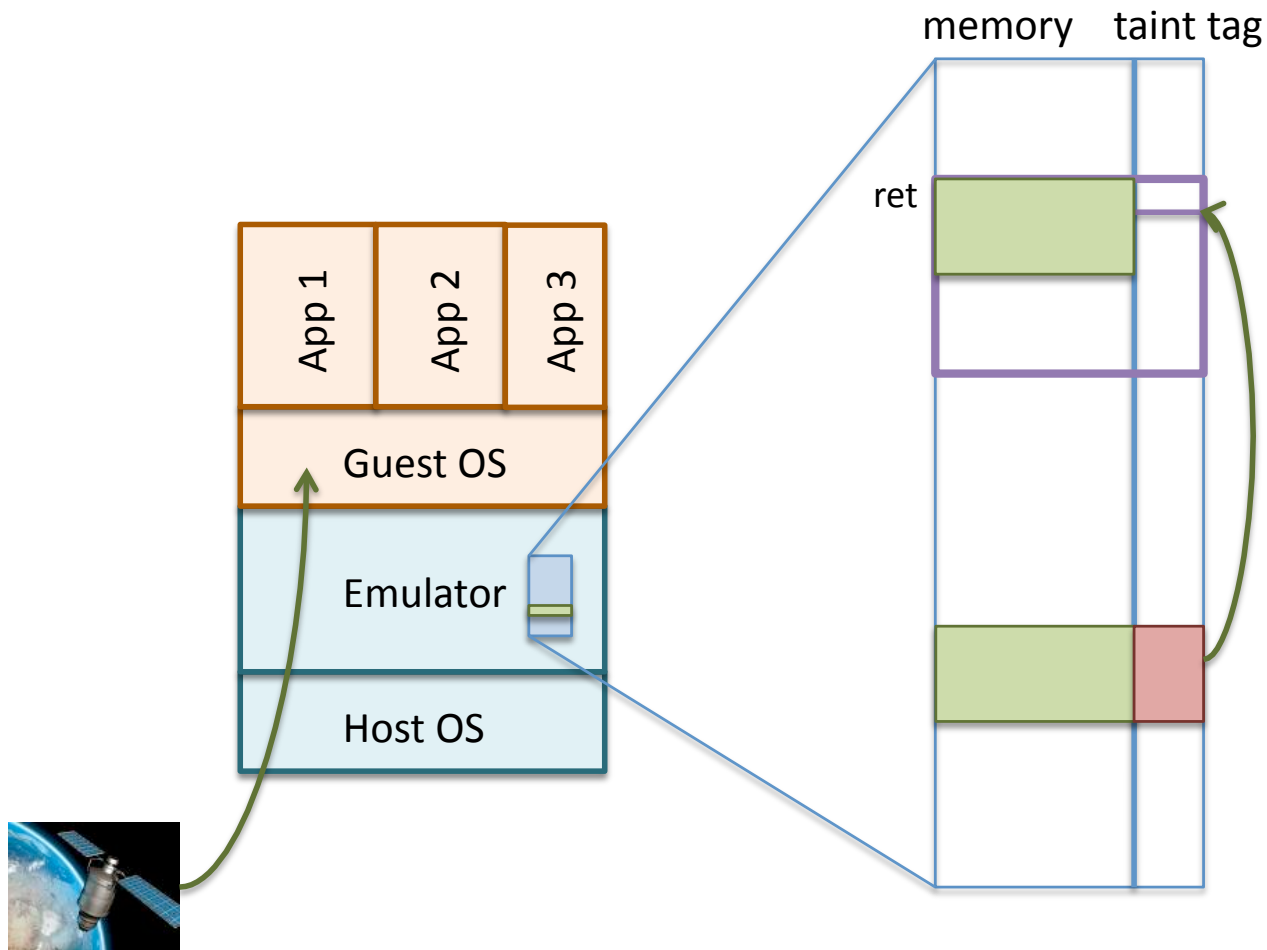


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2. Propagate taint through the OS.

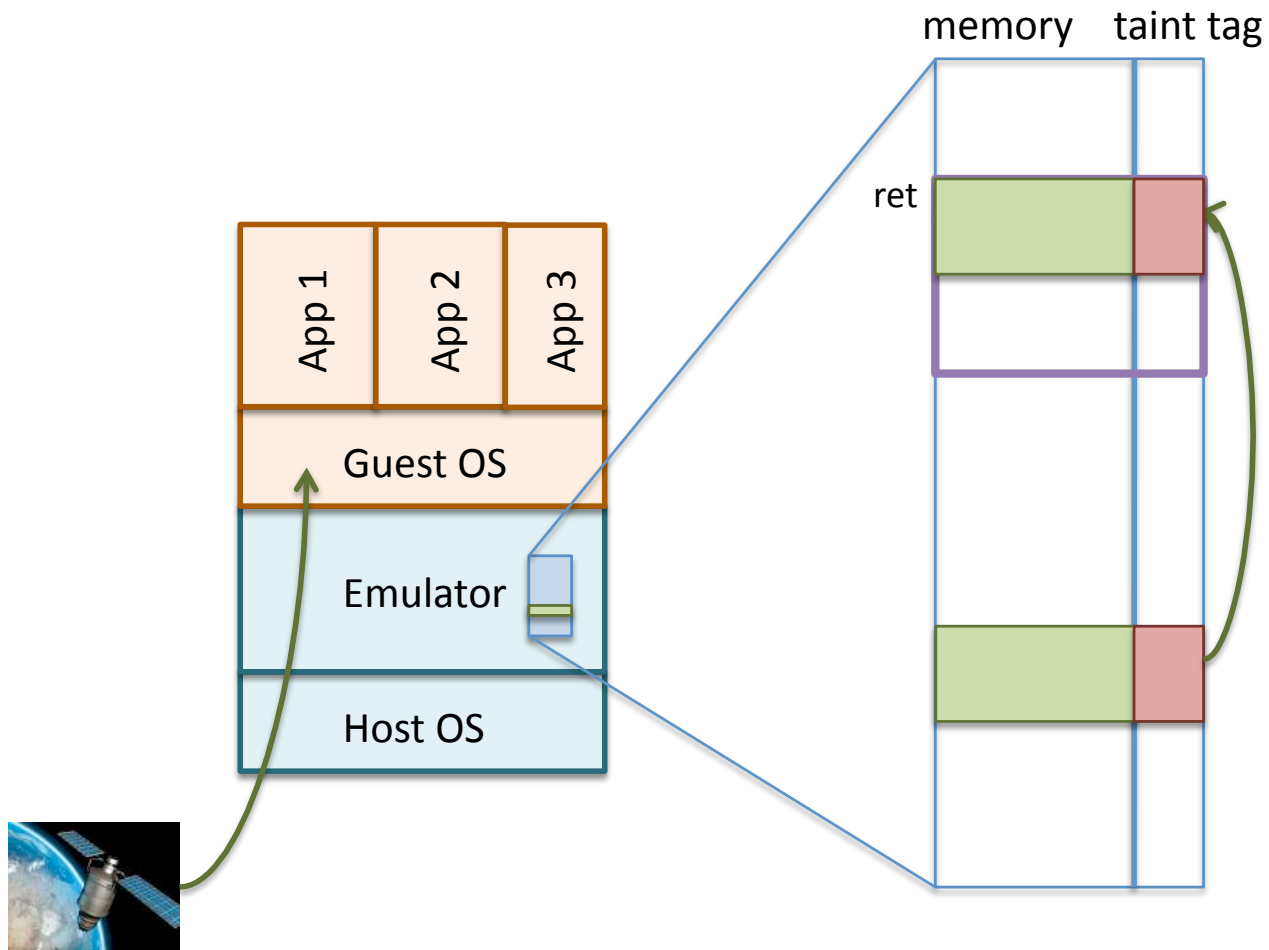


Basic tainting



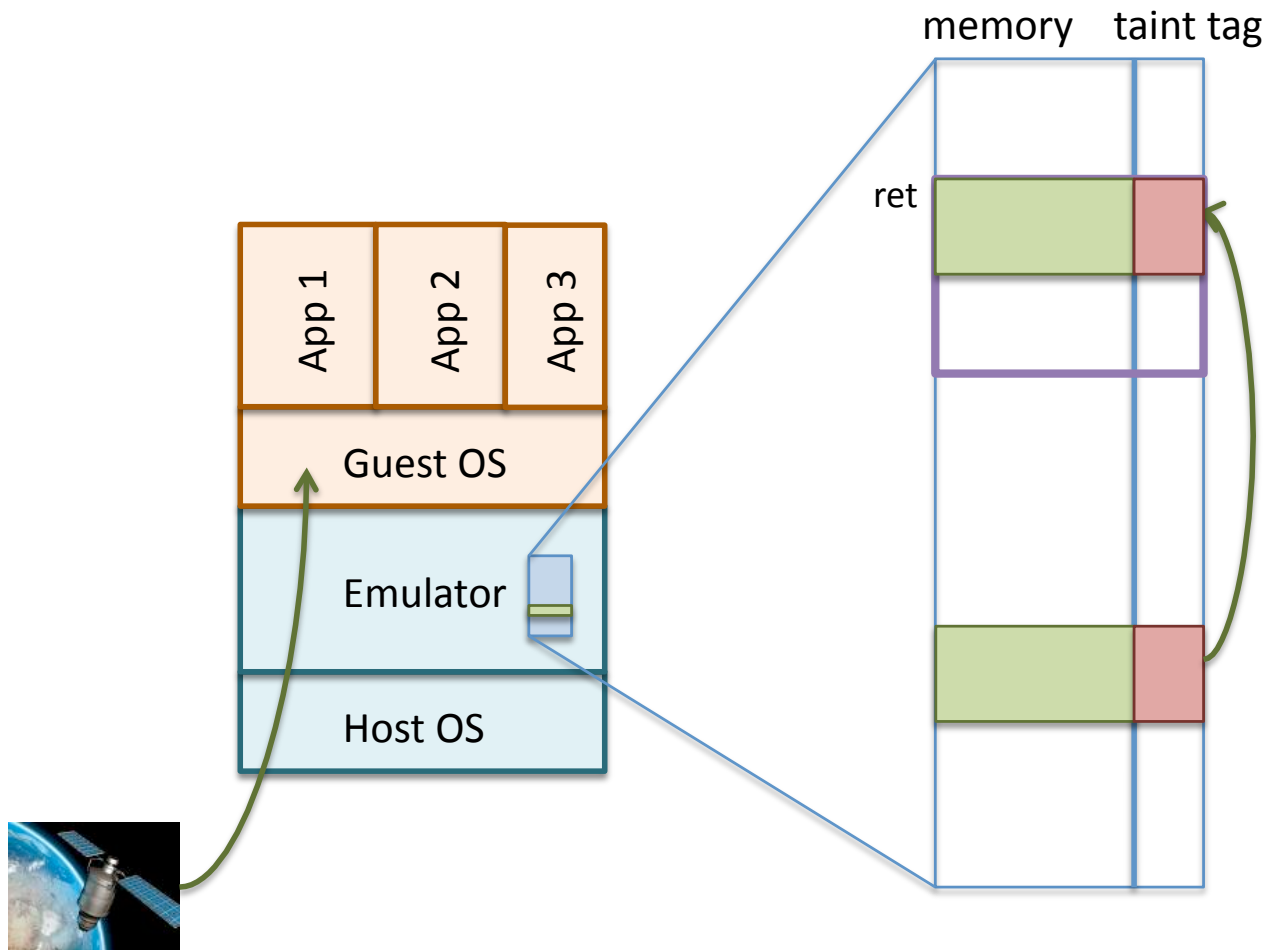
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Basic tainting



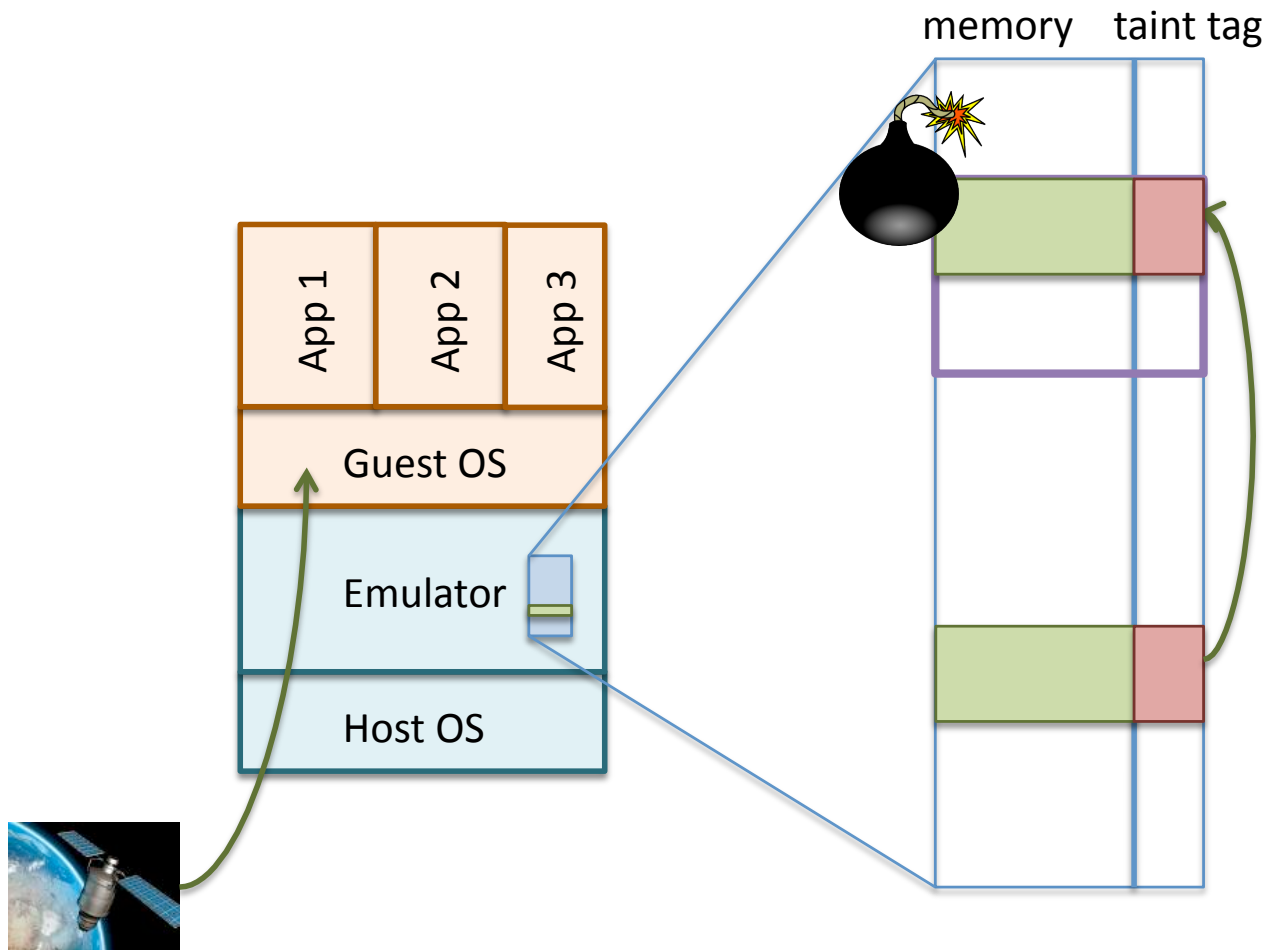
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Basic tainting



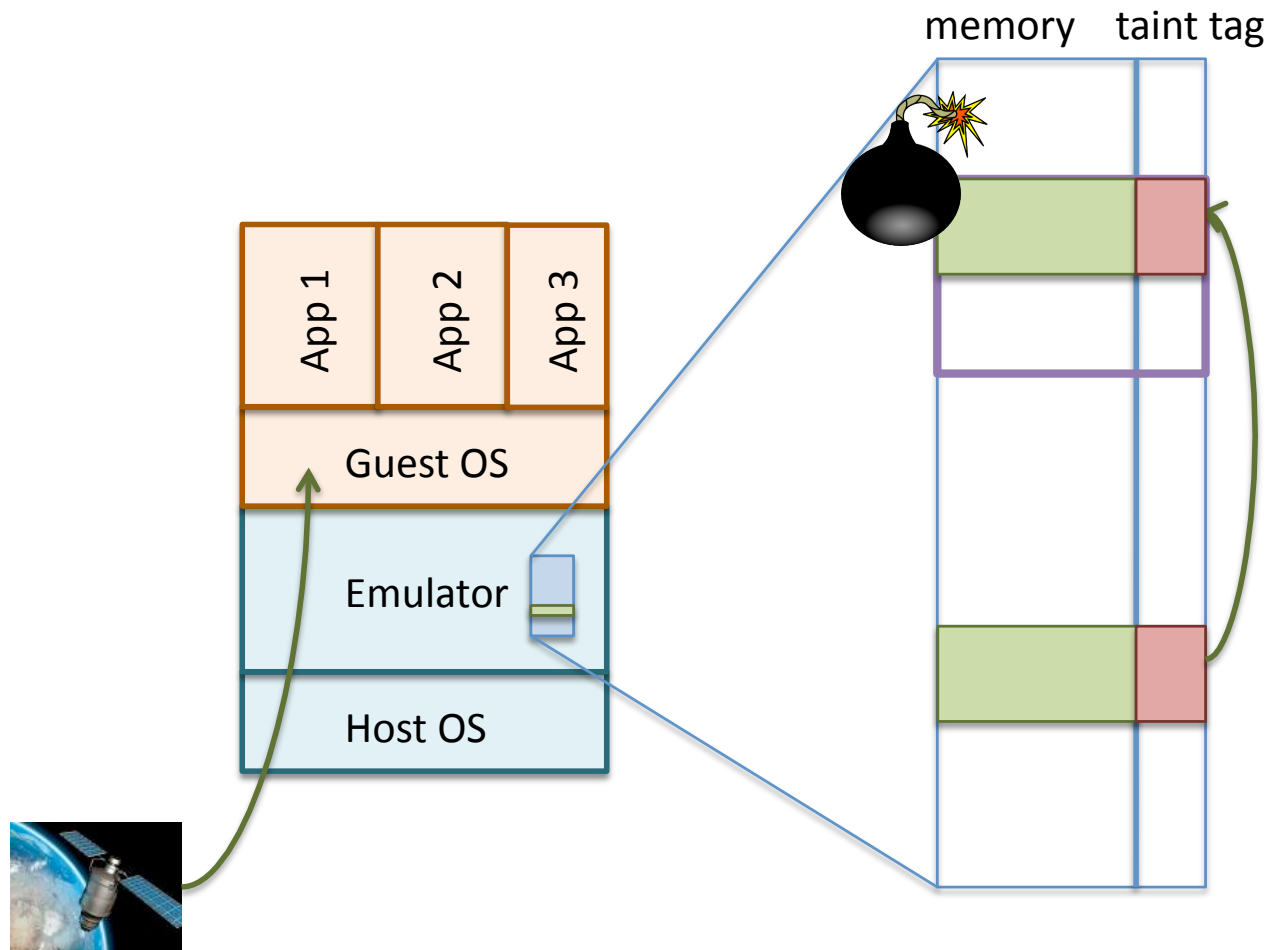
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3. Alert for dereferences due to tainted jumps, function calls/returns.

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Minos, MICRO 2004

Vigilante, SOSP 2005

Taintcheck, NDSS 2005

Argos, EuroSys 2006

Attacks: (in)effectiveness of basic tainting

```
void serve(int fd)
{
    char *reply = ...;
    char request[64];

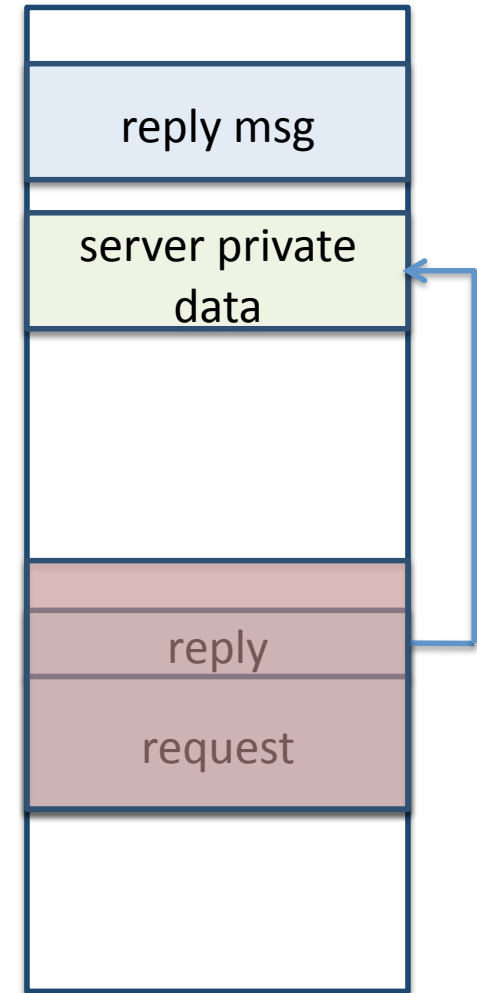
    read(fd, request, 128);
    srv_send(fd, reply, 1024);
}
```



Attacks: (in)effectiveness of basic tainting

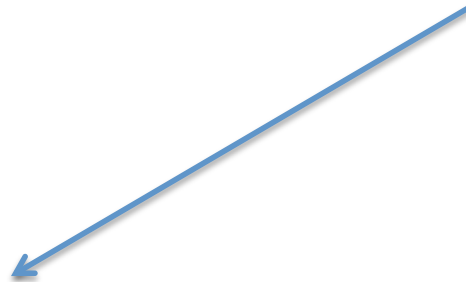
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Pointer tainting

1. Mark network data as tainted.
2. Propagate taint through the OS.



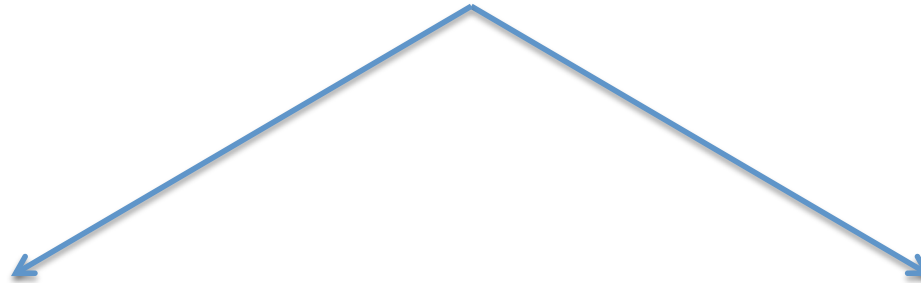
- Attacks

3. Alert for dereferences due to tainted jumps, function calls/returns.

+ If p is tainted, raise an alert on any dereference of p

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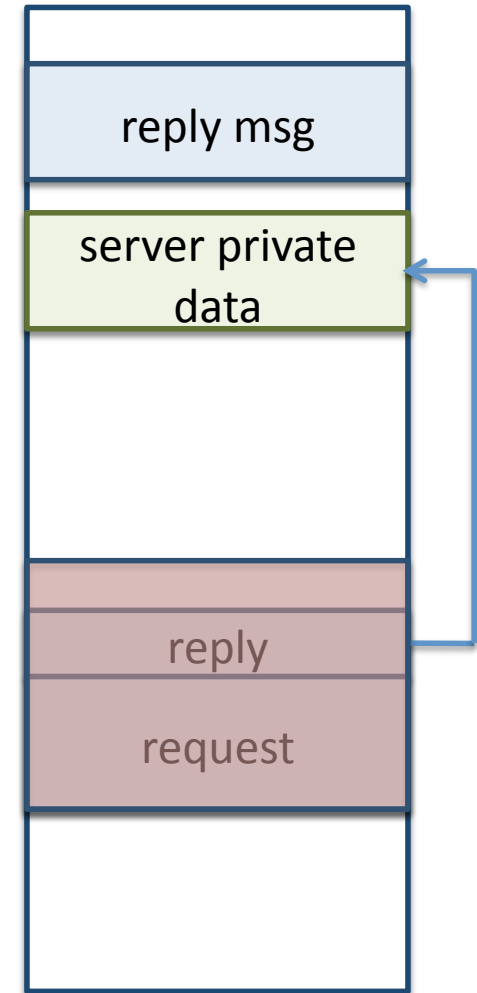
- Keylogger detection

+ If p is tainted, any dereference of p taints the destination

Attacks: effectiveness of pointer tainting

```
void serve(int fd)
{
    char *reply = ...;
    char request[64];

    read(fd, request, 128);
    srv_send(fd, reply, 1024);
}
```



Pointer tainting: FPs likely

```
void serve(int fd)
{
    char *reply;
    char request;

    read(fd, request, 1);

    srv_send(fd, reply, 1);
}
```

Pointer tainting: FPs likely

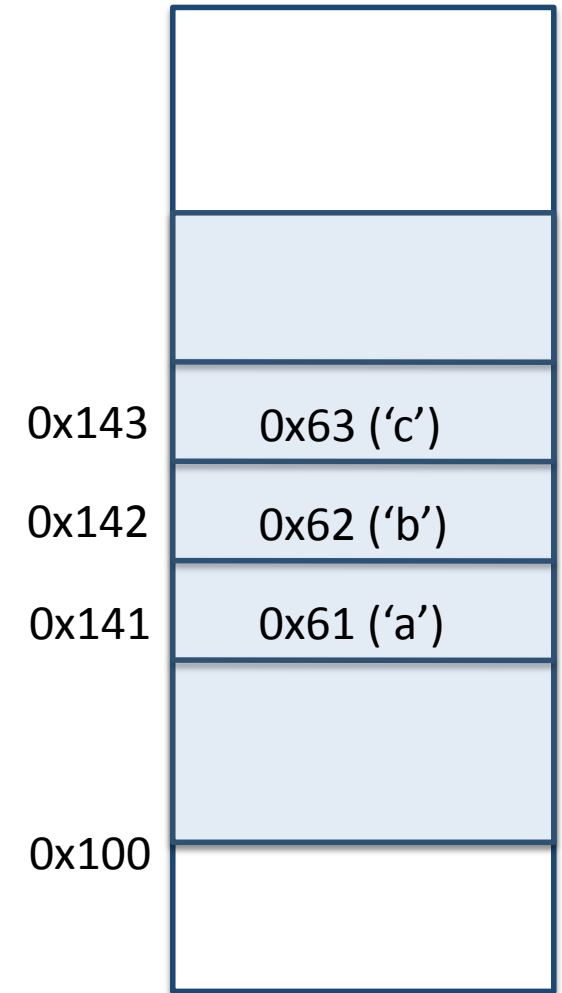
```
void serve(int fd)
{
    char *reply;
    char request;

    read(fd, request, 1);
    reply = to_lower[request];
    srv_send(fd, reply, 1);
}
```

Pointer tainting: FPs likely

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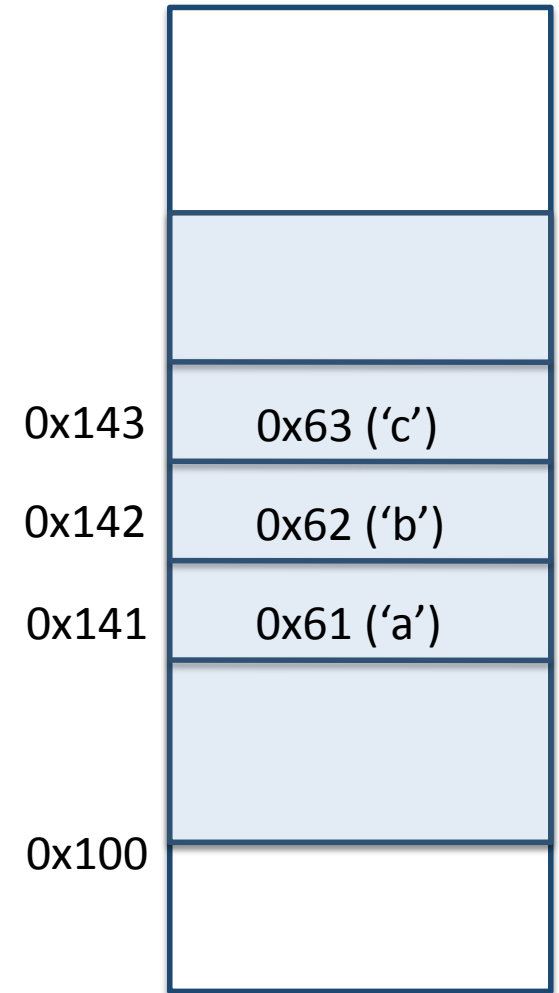


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    char *reply;
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    read(fd, request, 1);
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```

request = 0x41 'A'

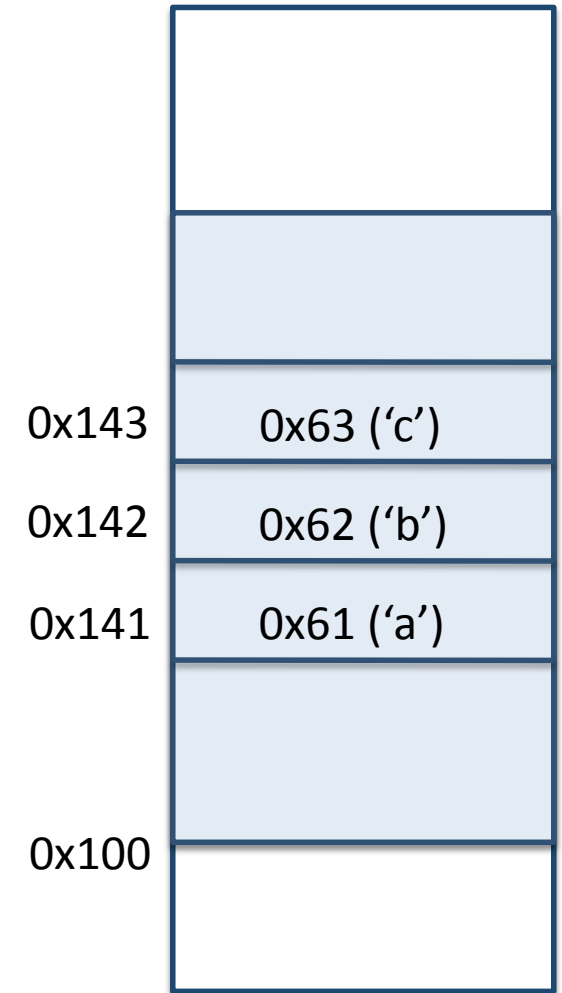


Pointer tainting: FPs likely

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void serve(int fd)
{
    char *reply;
    char request;

    read(fd, request, 1);
    reply = to_lower[request];
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}
```

```
request = 0x41  'A'
addr = 0x100 + request
reply = *addr
```



Pointer tainting: FPs likely

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void serve(int fd)
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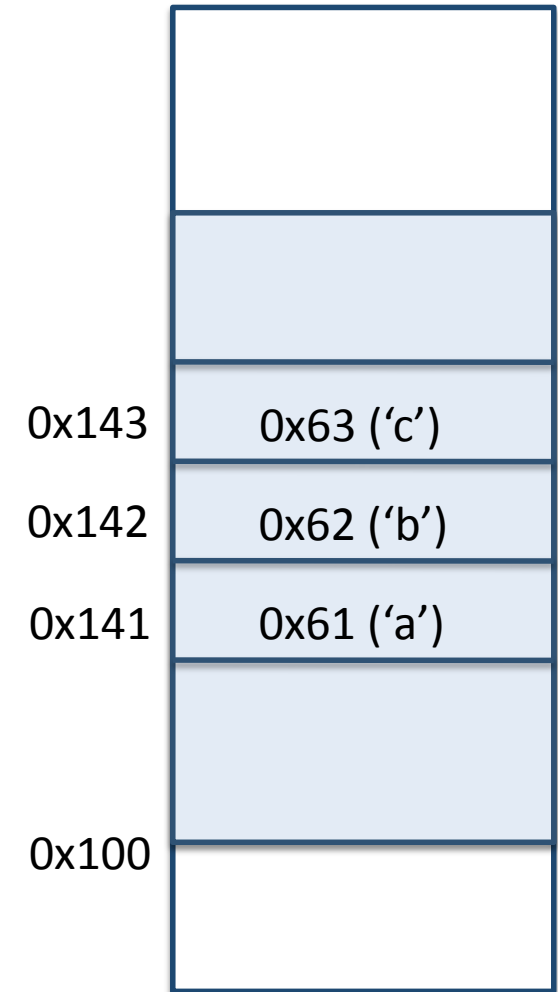
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`request = 0x41 'A'`

`addr = 0x100 + request`



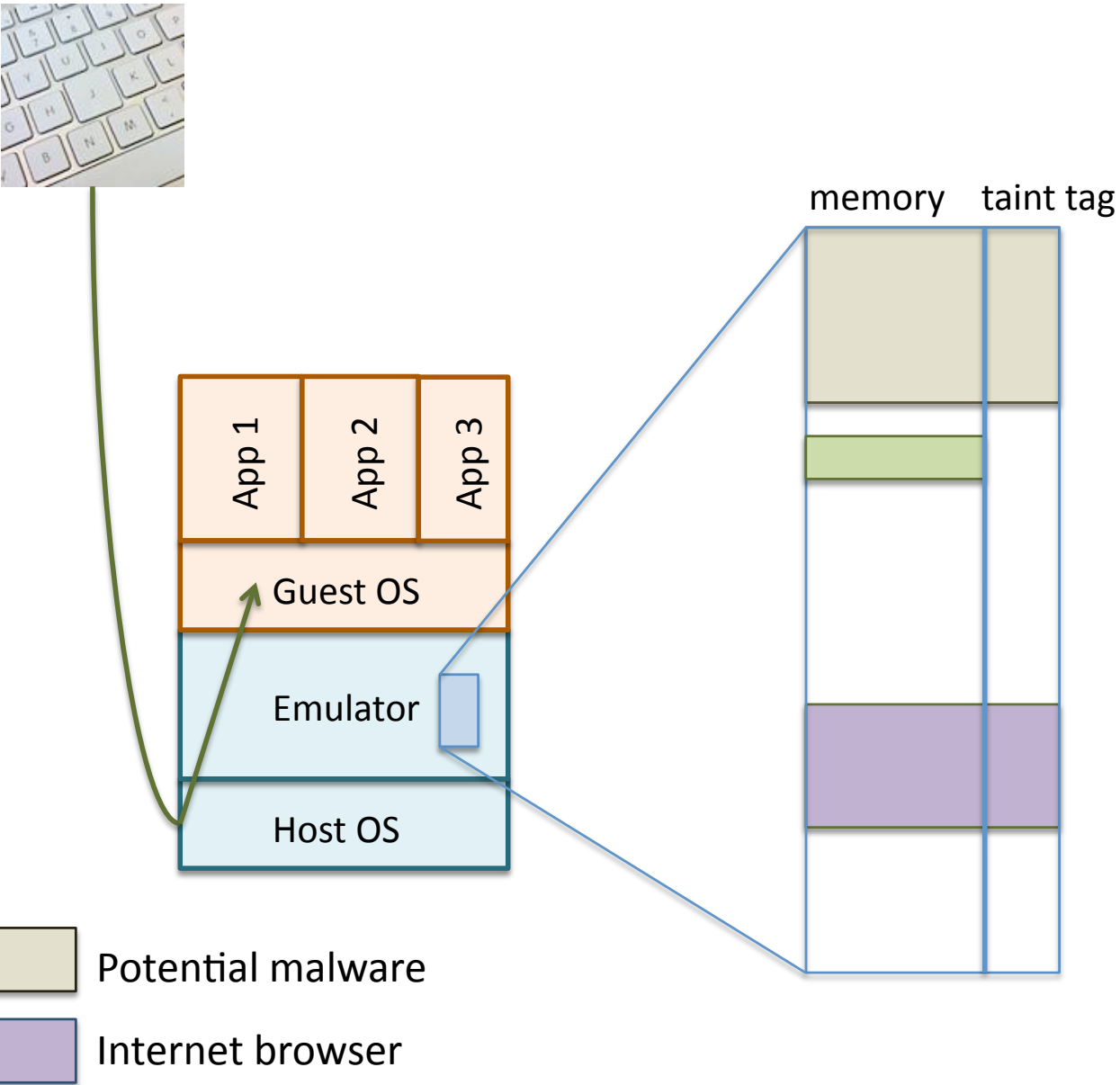
`reply = *addr`



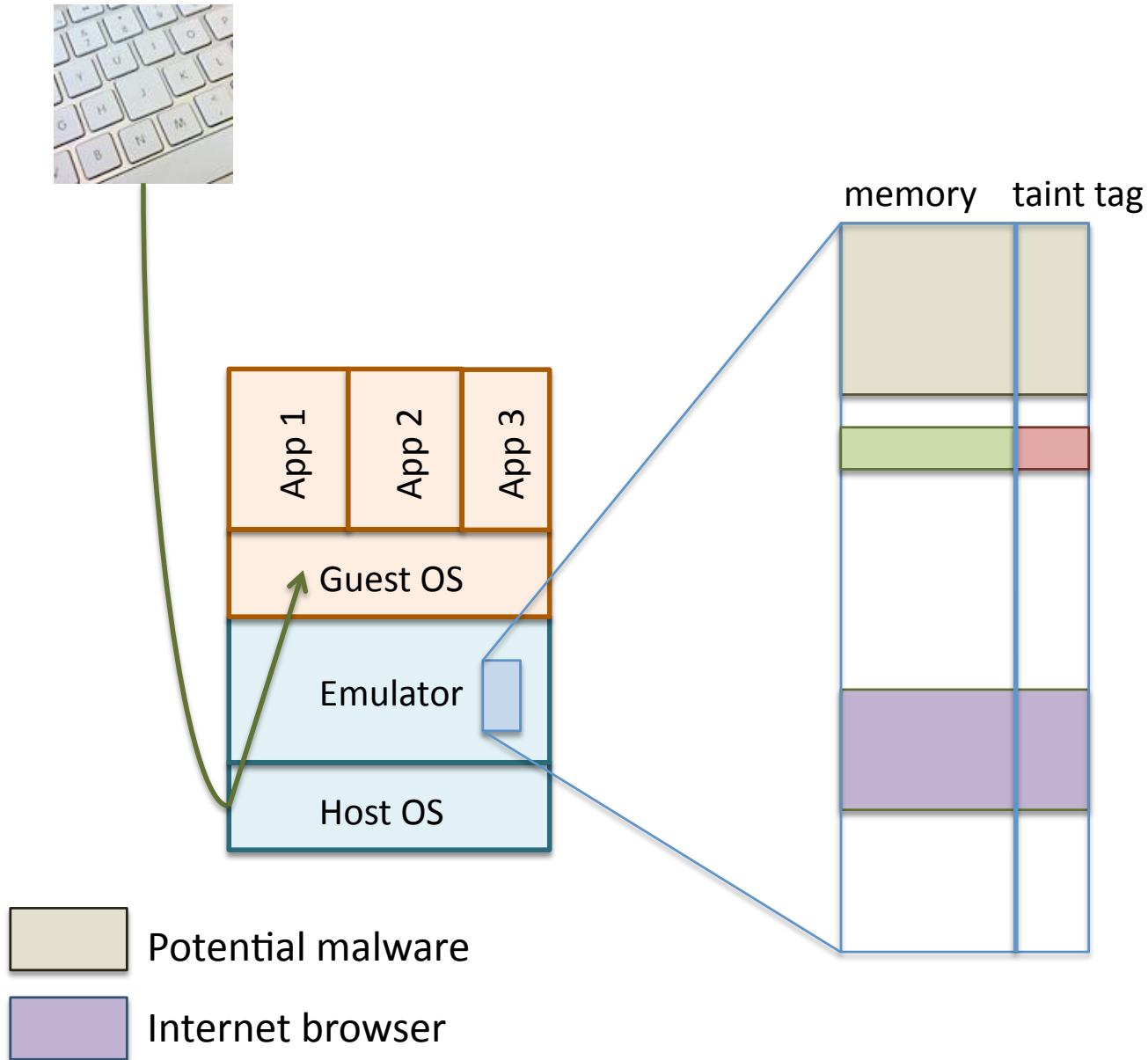
A close-up photograph of a white computer keyboard, showing keys for letters like Y, U, I, O, P, H, J, K, L, G, N, M, and B. The keys are slightly raised and have a clean, minimalist design.



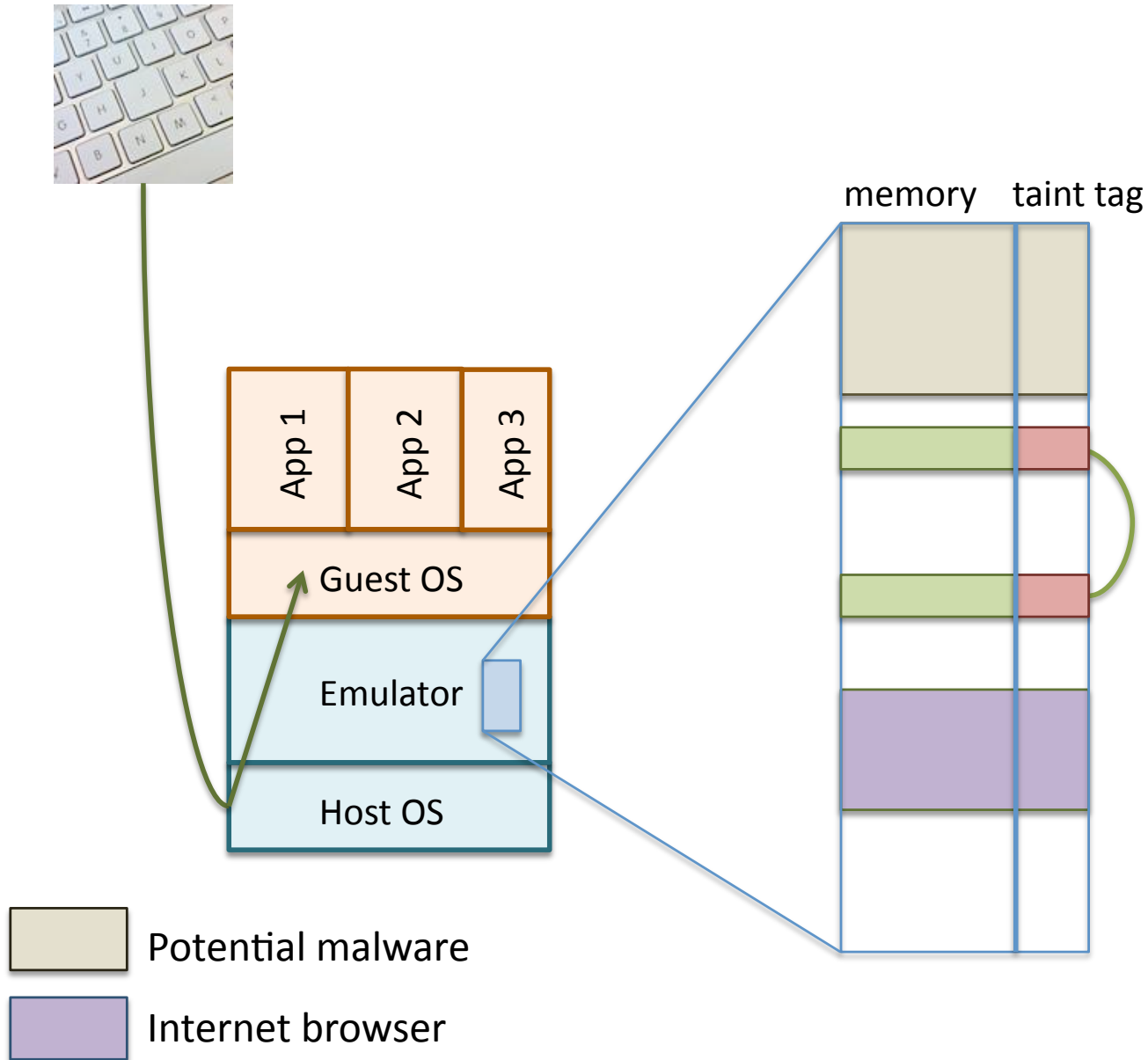
Keylogger detection



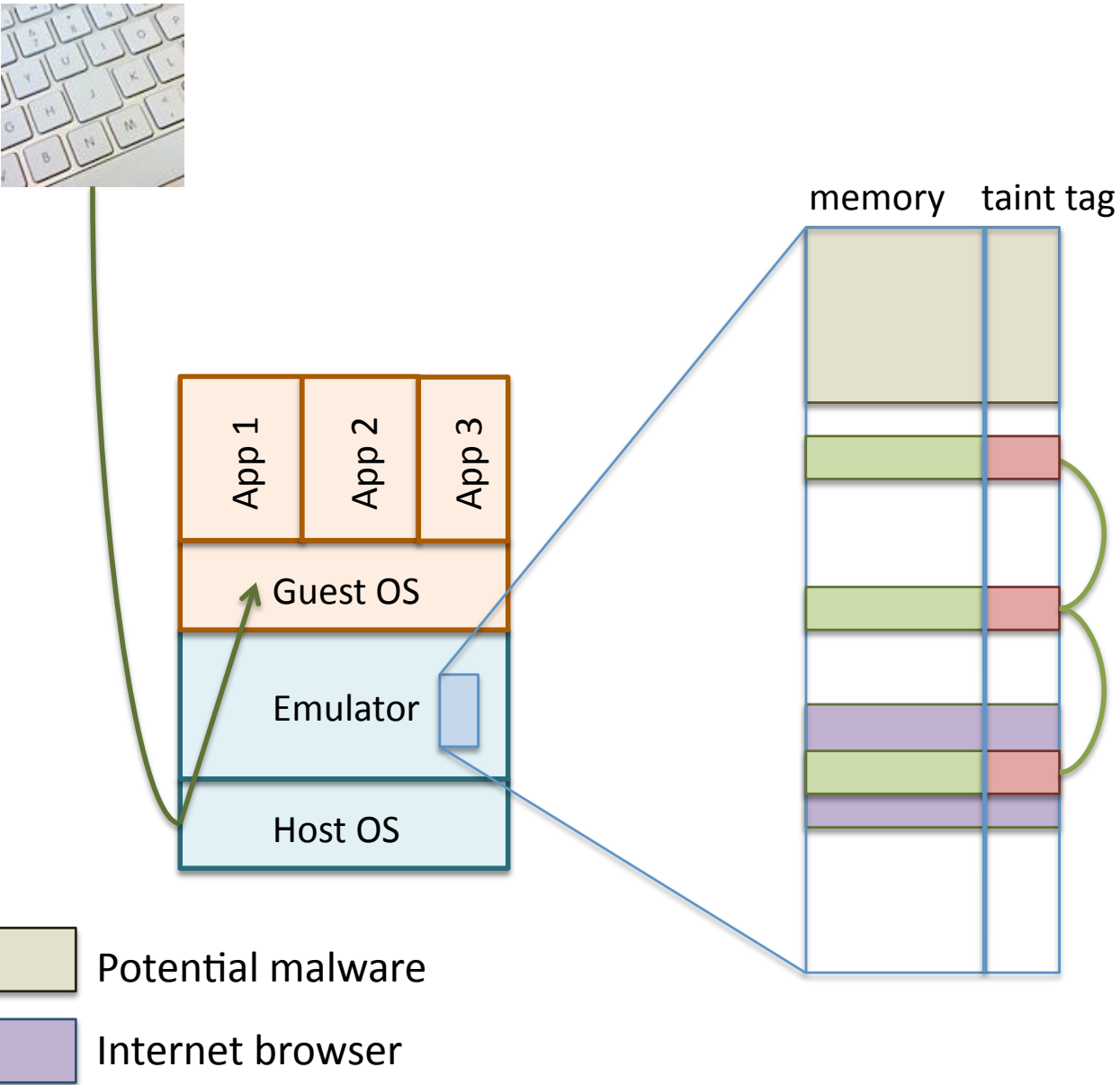
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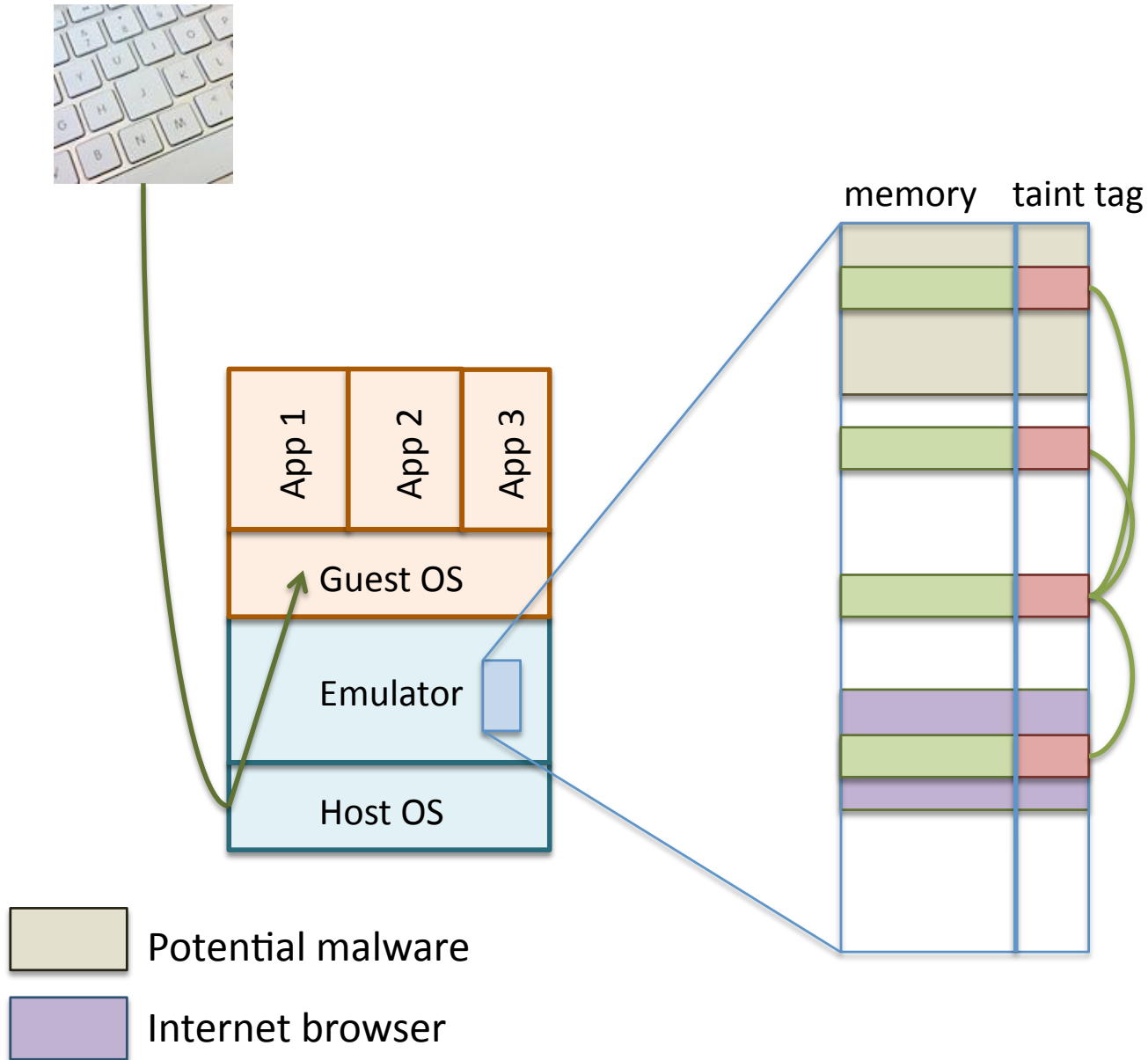
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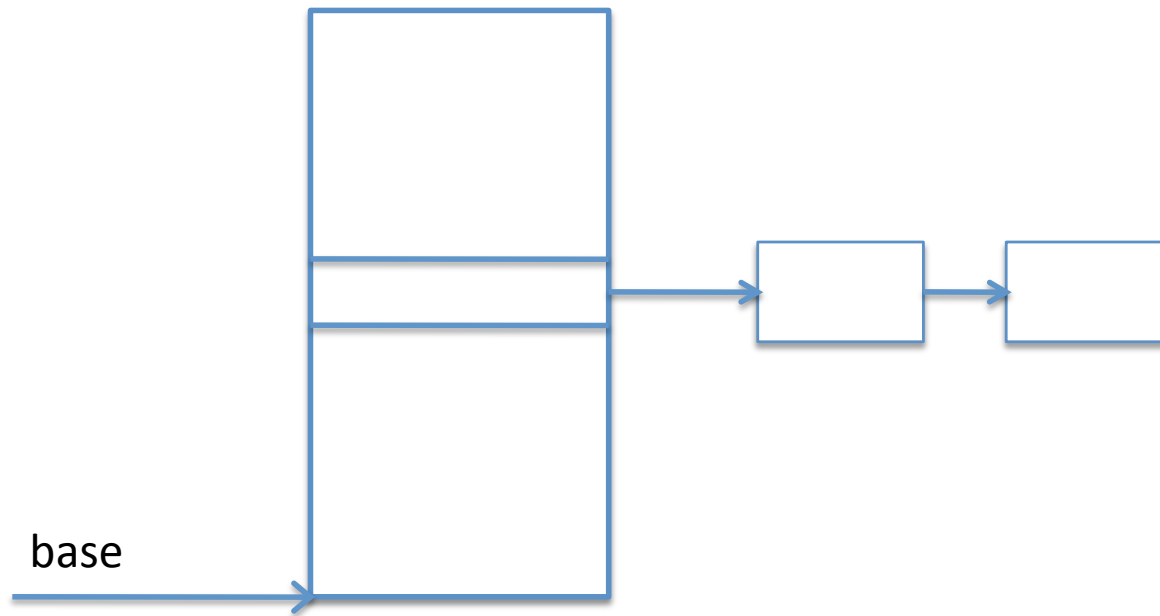


Keylogger detection: FPs likely (again)

```
struct hlist_head *head =  
    get_list_head(filename);
```

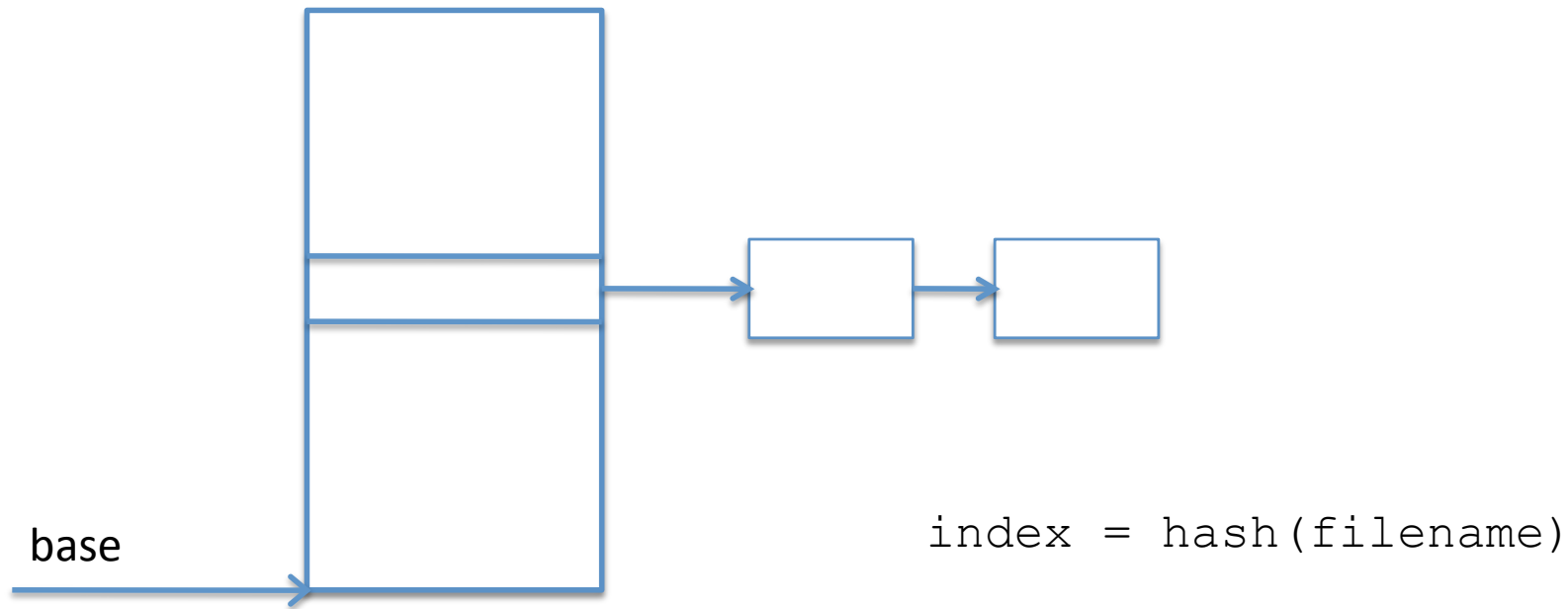
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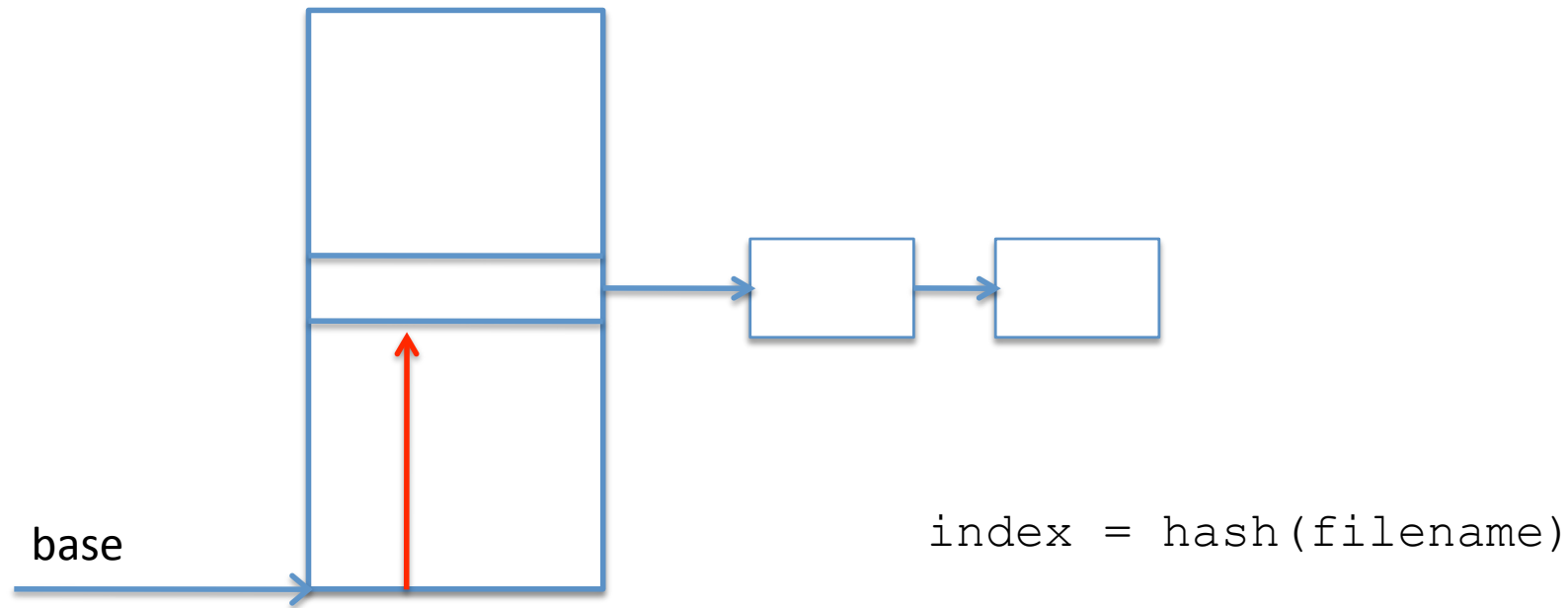
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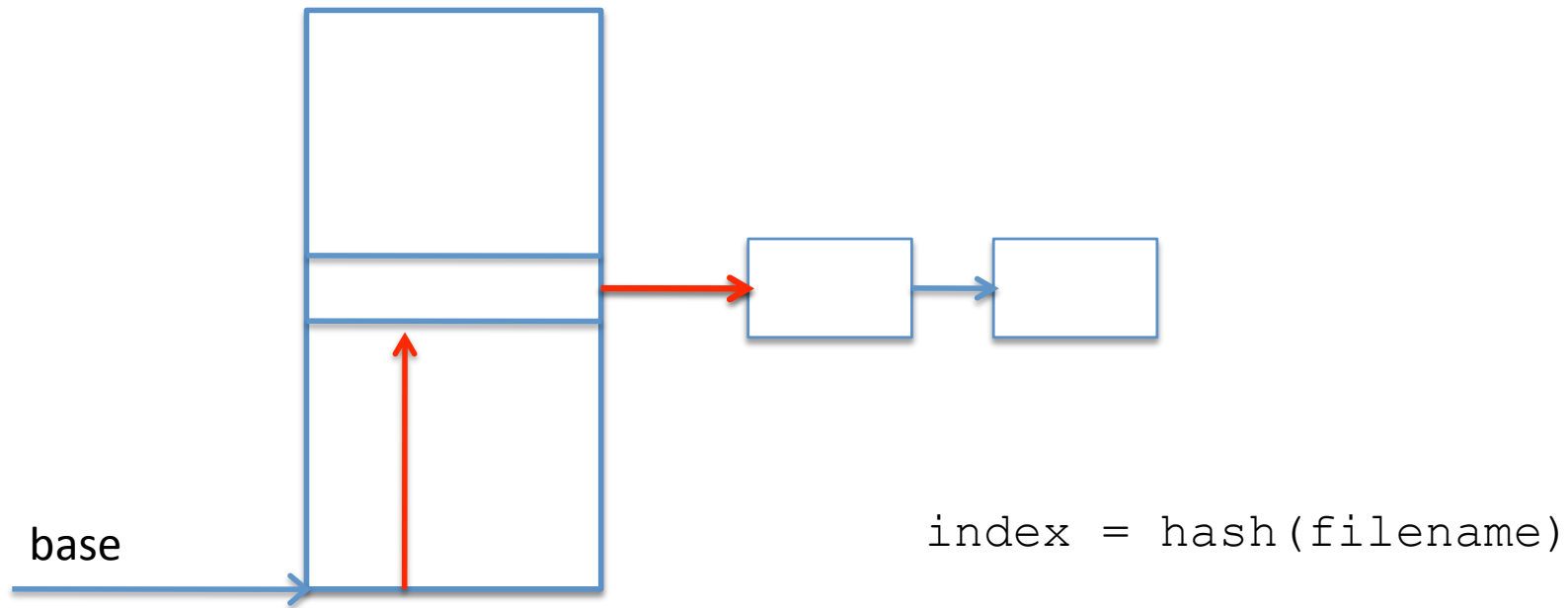
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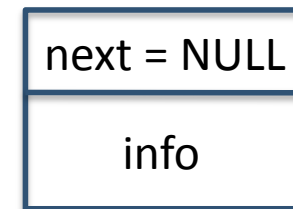
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struct dentry  
    *dentry = head->first;
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dentry of foo.txt

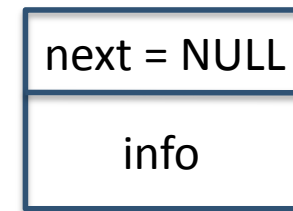


Keylogger detection: FPs likely (again)

```
struct hlist_head *head =  
    get_list_head("test.txt");
```

```
struct dentry  
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```

dentry of foo.txt

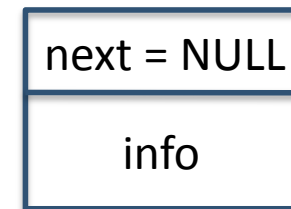


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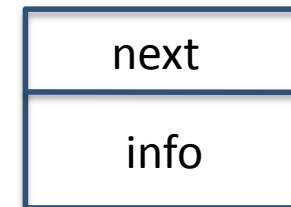
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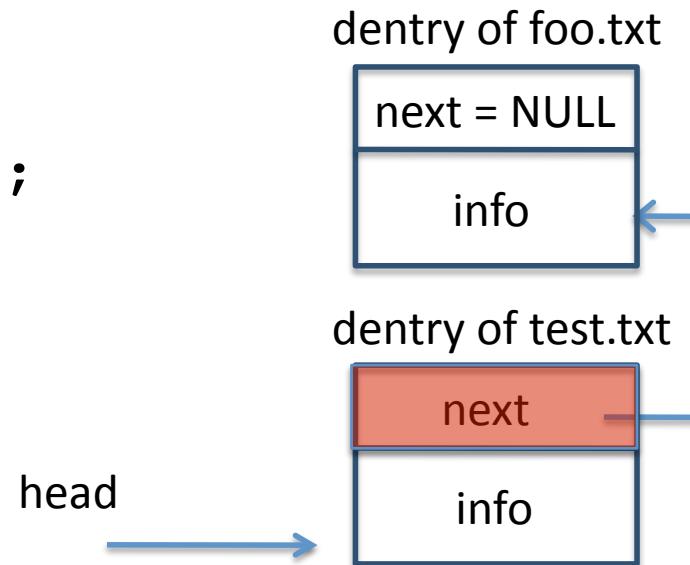
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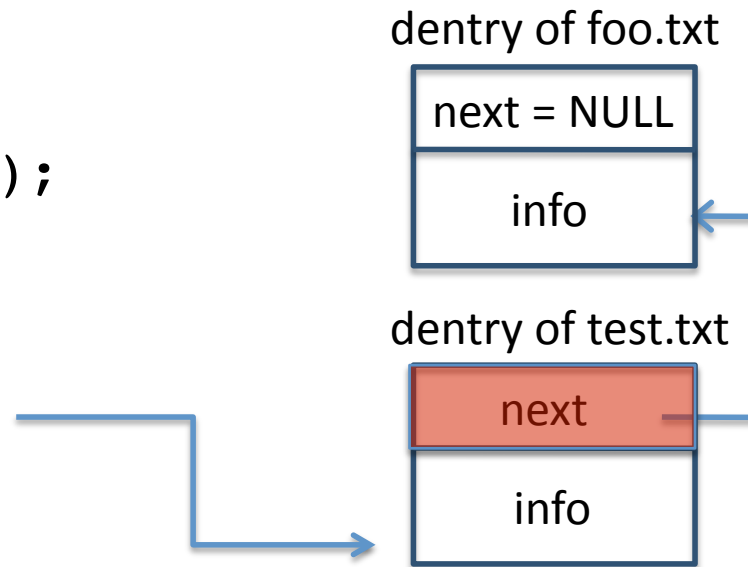
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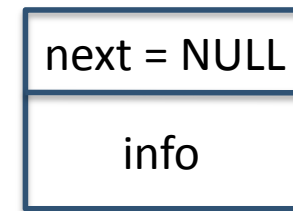
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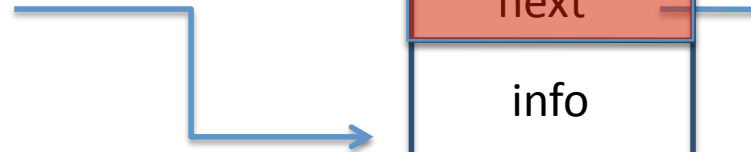
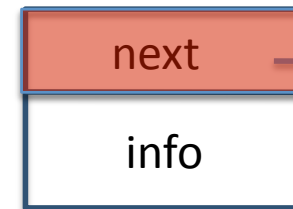
```
struct dentry  
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```
dentry = dentry->next;  
return info = dentry->info;
```

dentry of foo.txt



dentry of test.txt



Pointer tainting

1. Mark network data as tainted.
2. Propagate taint through the OS.



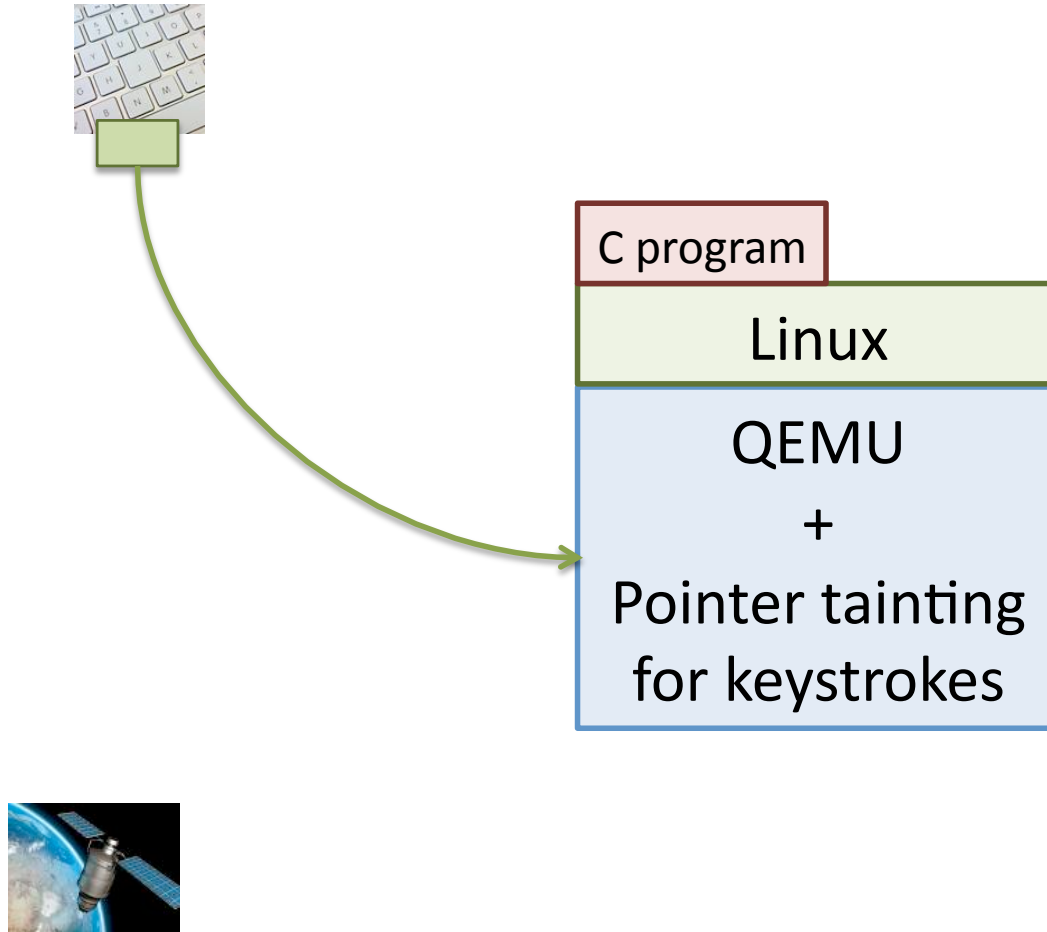
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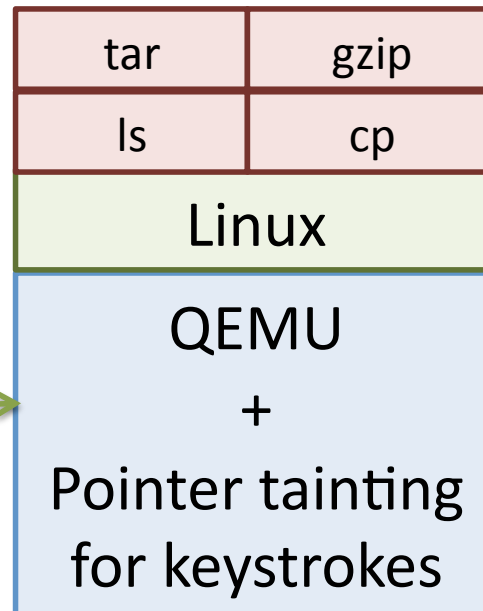
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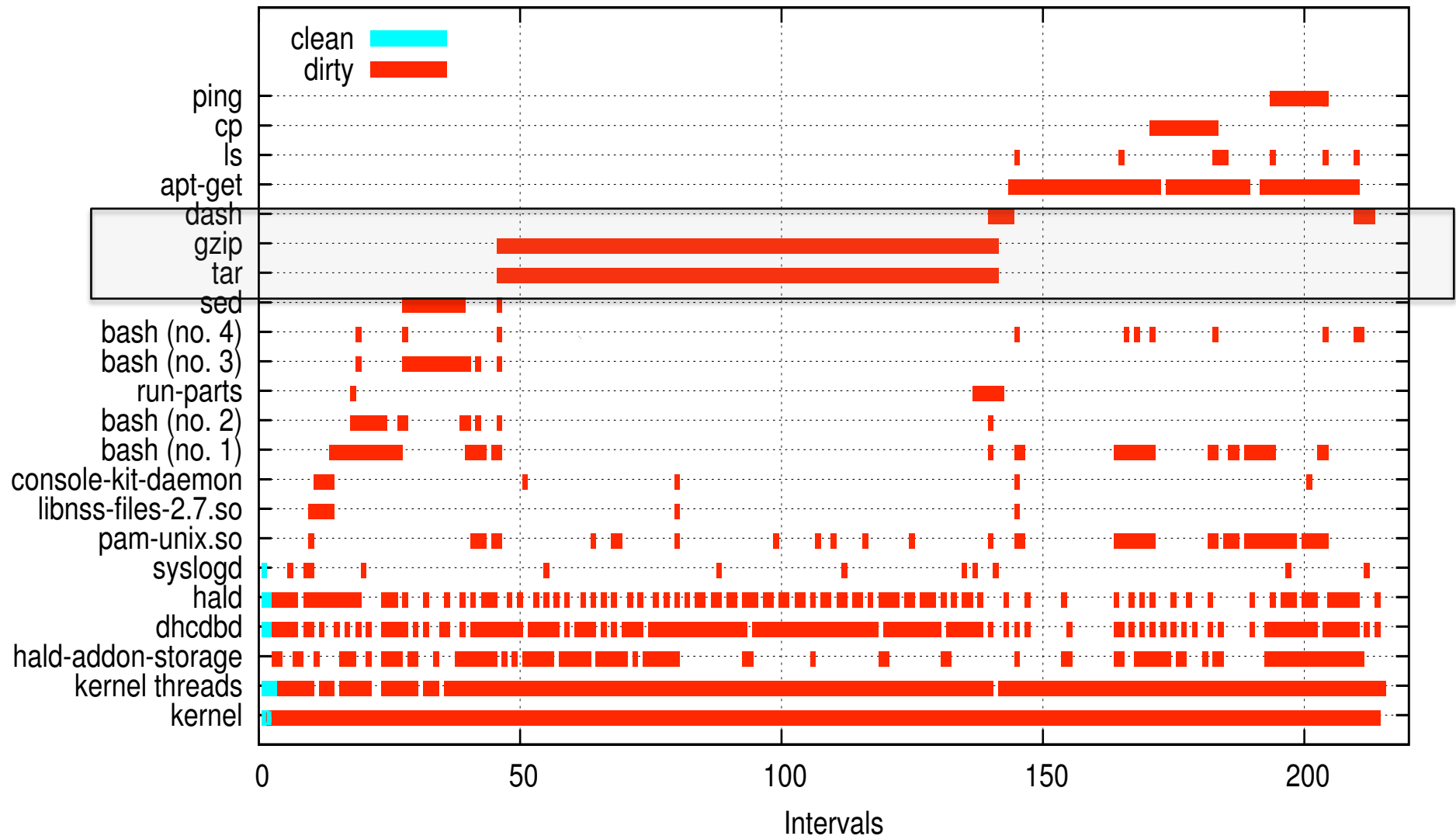
Experiment



Experiment



Keyloggers... false positives!



Containment

- White and black lists
- ESP/EBP protection
- Landmarking
- ...

Landmarking - motivation

Should NOT be tainted

```
struct dentry dentry =  
    prev_dentry->next;
```

Should be tainted

```
val = transl_table[index]
```

Landmarking - motivation

Should NOT be tainted

```
struct dentry dentry =  
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affect tainted address
with a **clean** value

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struct dentry dentry =  
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affect tainted address
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```
B = prev_dentry + offset
```

```
dentry = *B
```

Should be tainted

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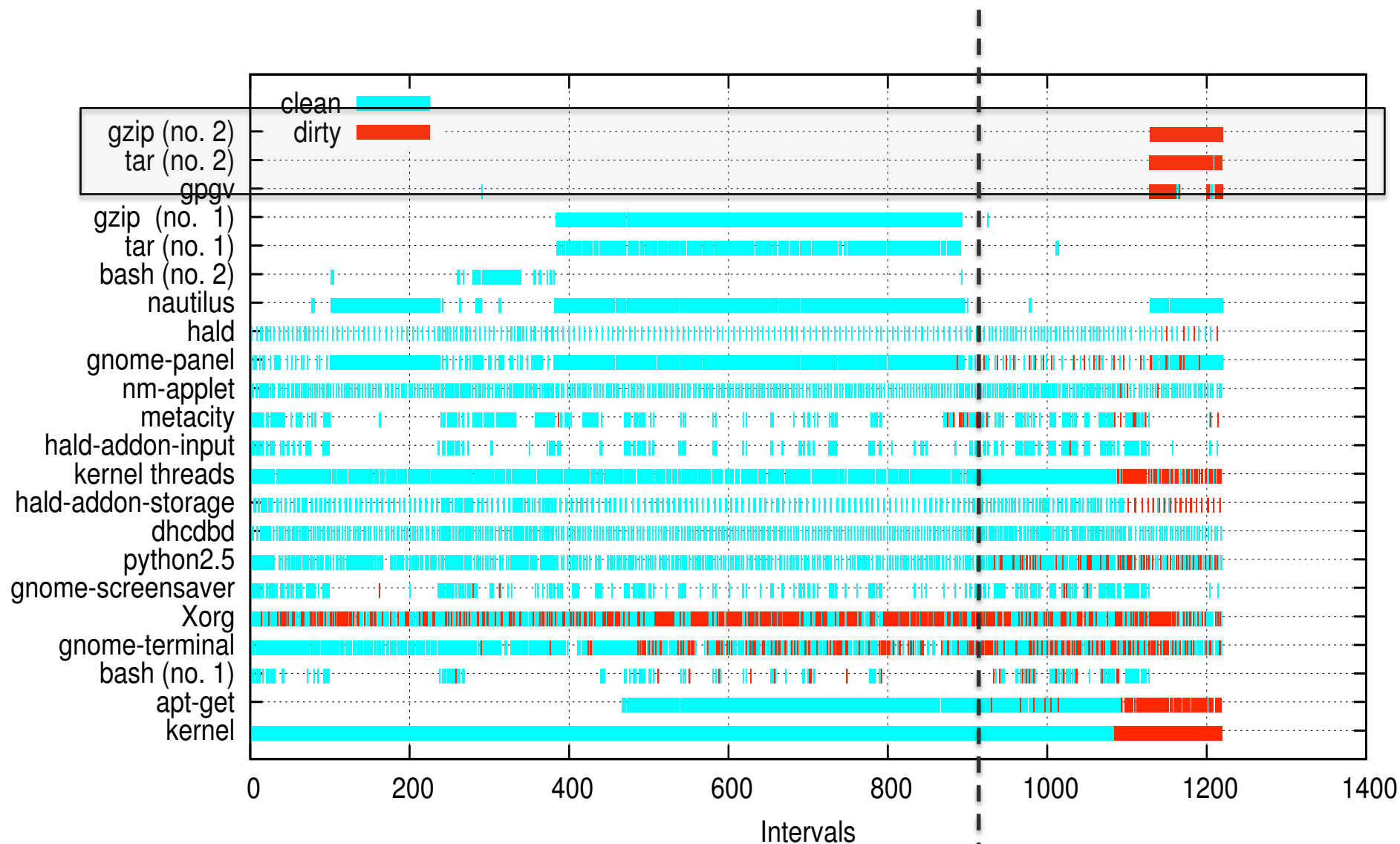
affect address
with a **tainted** value

A = address of an array

```
B = A + index*4
```

```
val = *B
```

Full containment - results



How bad are things?

Should NOT be tainted

```
struct hlist_head *head =  
    d_hash(parent, hash);
```

1. addr: combine clean pointer with a tainted index

```
struct dentry *dentry =  
    head->first;
```

- 2 new_addr: modify the resulting pointer with a constant

Should be tainted

```
attributes =  
    transl_table[kbd_data];
```

```
lower_case =  
    attributes->lower;
```

Conclusions

- We have analyzed pointer tainting
 - A popular technique for detecting memory corruption attacks and keyloggers
- Not suited for detecting privacy-breaching malware, like keyloggers
 - False positives hard to avoid
- Could be applied to detect memory corruption attacks
 - Not suitable for x86 and Windows

Backup slides

Pollution due to tainted ESP/EBP

- If ESP/EBP get tainted, taint spreads instantly
 - `mov eax, dword ptr [ebp + 08h]`
 - `pop eax`
- How ESP/EBP can become tainted?
 - Linux kernel has numerous places where it can happen,
 - E.g., a common operation like opening a file ends up tainting EBP,
 - Details in the paper

Pollution due to pointer arithmetic

Should NOT be tainted

```
struct fd {  
    HANDLER handler;  
    STRING filename;  
    struct fd *next;  
};
```

A = address of filename

B = **A** - 0x0004

h1 = *****(**B** + 0x0000)

fd2 = *****(**B** + 0x0020)

Should be tainted

A = address of an array

i = index to be accessed

B = A + **i***4

Translated value:

val = ***B**

Pollution due to pointer arithmetic

Should NOT be tainted

```
struct fd {
```

```
    HANDLER
```

```
    STRING filename;
```

```
    struct fd *next;
```

```
};
```

A = address of filename

B = **A** - 0x0004

h1 = *****(**B** + 0x0000)

fd2 = *****(**B** + 0x0020)

Should be tainted

A = address of an array

i = index to be accessed

B = A + **i***4

Translated value:

val = ***B**

How to distinguish between
these two cases?

Landmarking

```
typedef struct test_t {  
    int i;  
    struct test_t *next;  
} test_t, *ptest_t;  
ptest_t table[256] = ...;
```

```
ptest_t i1 = table[index];    // tainted  
    A = (table+index*sizeof(test_t))
```

```
ptest_t i2 = i1->next;        // clean  
    addr: *(A + offset(next))
```

```
int i3 = i1->i;                // clean  
    addr: *(A)
```

Landmarking – why FPs?

- Possible scenarios:
 - Assume `eax` contains a calculated tainted address
 - It can be copied and altered before dereference
 - Then both values become tainted
 - Addresses calculated directly
 - an array `A` of struct `{int a; int b;}`
 - `A[index].b: int b = *((char*)A+8*index+4)`
 - Very simplistic, but the same problem might hold for queues, stacks and hashtables

Landmarking – more problems

- False negatives
 - Translation table containing structures instead of single elements

```
attributes = transl_table[kbd_data];  
lower_case = attributes->lower;
```

- Much more problems in the paper